

# REALTINE OPERATING SYSTEMS

REALTIME OPERATING System is said to be Real Time if it is required to complete it's work & deliver it's services on time. Example – Flight Control System All tasks in that system must execute on time. Non Example – PC system

# HARD AND SOFT REAL Hard Refer TSNES/Ster MS

Validation by provably correct procedures or extensive simulation that the system always meets the timings constraints Failure to meet deadlines is fatal example : Flight Control System

## SOFT REALTIME SYSTEM

- Demonstration of jobs meeting some statistical constraints suffices.
  - Late completion of jobs is undesirable but not fatal.
  - System performance degrades as more & more jobs miss deadlines
  - Online Databases

ROLE OF AN OS IN BEAL THE STENS Often no OS involved Micro controller based Embedded Systems

#### Some Real Time Applications are huge & complex

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<sup>D</sup>Multiple threads

Complicated Synchronization Requirements Filesystem / Network / Windowing support OS primitives reduce the software design time



\* Resource Allocation.

Interrupt Handling.

Other issues like kernel size.

# <u>SCHEDULING IN</u> <u>RIOS</u>

More information about the tasks are known No of tasks Resource Requirements Release Time Execution time Deadlines

Being a more deterministic system better scheduling algorithms can be devised.

#### SCHEDULING ALGARICHIMS IN RTOS Weighted Round Robin Scheduling

Priority Scheduling (Greedy / List / Event Driven) CLOCK DRVEN All parameters about jobs (release time/ execution time/deadline) known in advance. Schedule can be computed offline or at some regular time instances. Minimal runtime overhead. Not suitable for many applications.

### MEIGHTED ROUND ROBIN

- Jobs scheduled in FIFO manner
- Time quantum given to jobs is proportional to it's weight
- Example use : High speed switching network QOS guarantee.
- Not suitable for precedence constrained jobs. Job A can run only after Job B. No point in giving time quantum to Job B before Job A.

#### PRIORITY SCHEDULING

(Greedy/List/Event Driven) Processor never left idle when there are ready tasks Processor allocated to processes according to priorities Priorities

static - at design time

Dynamic - at runtime

PRORTY Easiest beachelerst Ceder) Process with earliest deadline given highest priority Least Slack Time First (LSF) slack = relative deadline - execution left Rate Monotonic Scheduling (RMS) For periodic tasks Tasks priority inversely proportional to it's period

RESOURCE RALL-ORIATION IN The issues with scheduling applicable here. Resources can be allocated in Weighted Round Robin Priority Based

Some resources are non preemptible Example : semaphores Priority Inversion if priority scheduling is used

#### INTERRUPT Internel and the same an isocritical sections of the kernel No worst case bound on interrupt latency available g: Disk Drivers may disable interrupt for few hundred milliseconds Not suitable for Real Time Applications Interrupts may be missed

TWOLEVELNDERRUPT TRANSPICE Handling Top Half Interrupt Handler Called Immediately – Kernel never disables interrupts Cannot invoke thread library functions - Race Conditions Bottom Half Interrupt Handler Invoked when kernel not in Critical Section Can invoke thread library functions Very Low Response time (as compared to Linux)

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# <u>OTHERFEATURES</u>

Footprint Small footprint (~50kb)

Oskit's Device Driver Framework Allows direct porting of existing drivers from Linux. Example – Ethernet Driver of Linux